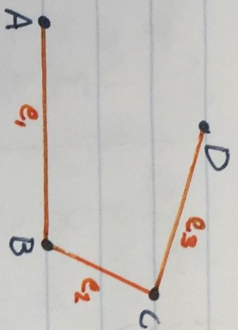


## Chapter 10: Graphs (Deviates from textbook)

### Graphs

\* Graph (non-directed) example:



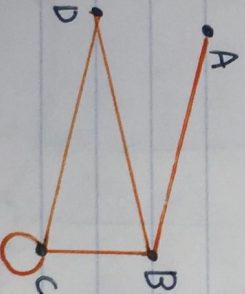
Vertices =  $\{A, B, C, D\}$   
 Edges =  $\{AB, BC, CD\}$

\* Adjacent = If two vertices are connected by an edge, they are said to be **adjacent** to each other.

- An edge that is incident on a vertex means that the edge is attached to that vertex.

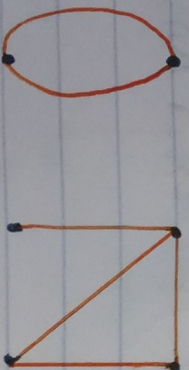
\* Types of Graphs:

1.) Connected Graph



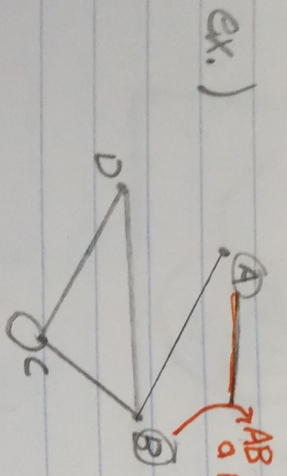
\* Bridge = An edge whose removal turns a connected graph into a disconnected graph

2.) Disconnected Graph



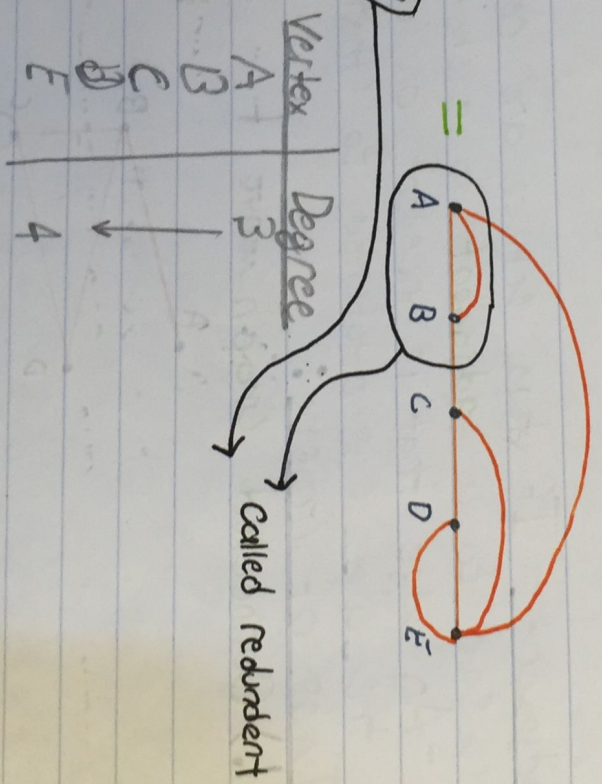
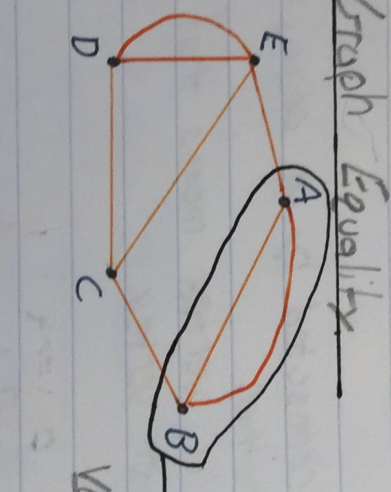
this is the whole graph

\* Degree of a Vertex: The degree of a vertex = to the # of edges that are incident on the vertex.  
 in other words, loops count twice!



Vertex	Degree
A	1
B	3
C	4
D	2

\* Graph Equality



3.) Complete Graph = A graph where each vertex is adjacent to every other vertex ( $K_n$ )

ex.  $K_4$

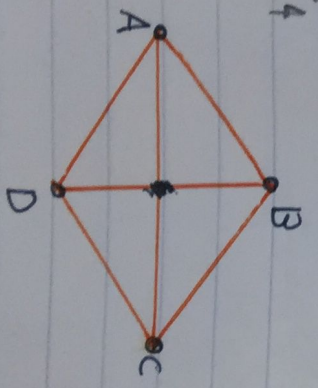


diagram created with draw.io

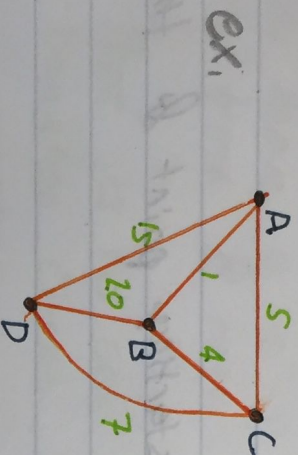
4.) Weighted Graphs: The edge values will have numerical values that will be the "weight" of the edge.

What does this "weight" represent?  
Depends on the problem!

- Miles / distance

- Time

- Cost

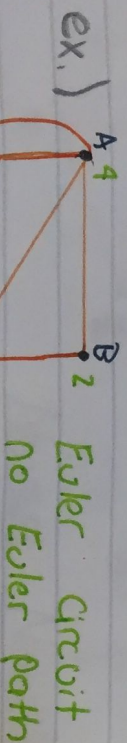
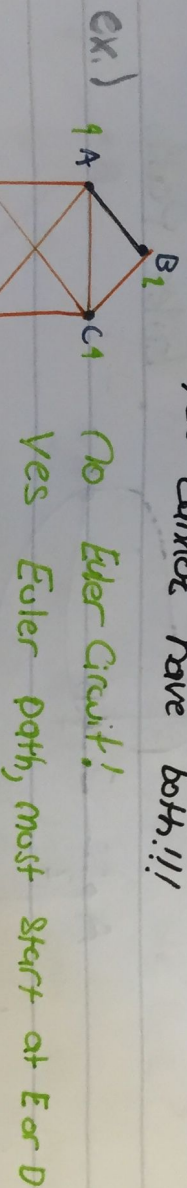


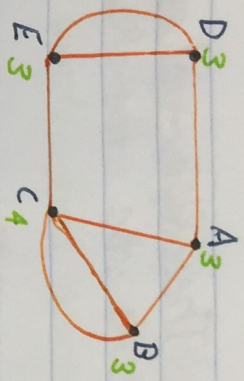
### Circuits / Paths (Euler)

\* Euler Circuit: You leave from a vertex, travel each edge exactly once and return to where you started.

Euler Path: You leave from a vertex, travel each edge exactly once and end up somewhere else.

You cannot have both!!!





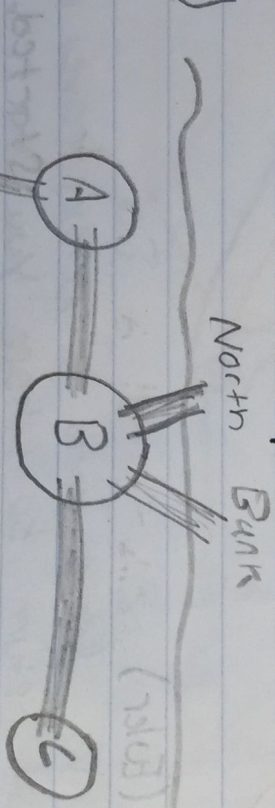
Neither Euler Circuit  
or Euler  
Path!

Look @ the degrees of these graphs !!!

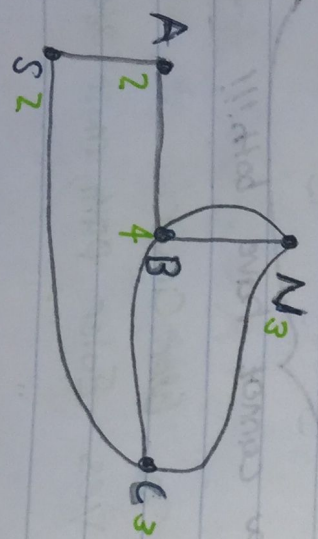
- To have an Euler Circuit a graph must be connected and each vertex must have an even degree
- To have an Euler Path a graph must be connected and exactly two of the vertices must have an odd degree. (others even)

↓  
one odd degree will be the starting point & the other the end point

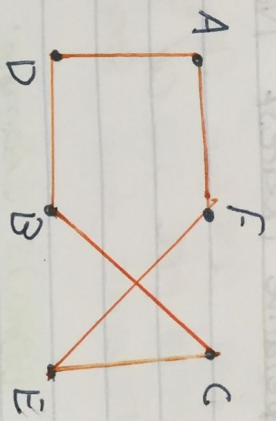
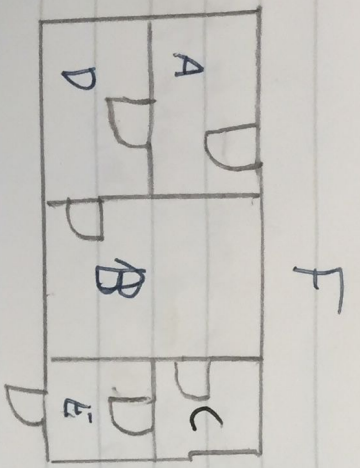
ex. 1)



North Bank



Euler Path



Euler Circuit

Hamilton Circuits

\* You leave a vertex, travel to every other vertex exactly one time & return to the starting vertex.

There is no easy way to determine if a graph has a hamilton circuit.

Traveling Salesperson Problem

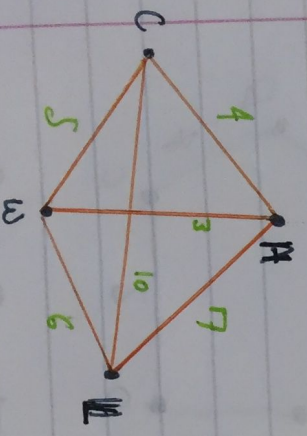
\* Find the most efficient way to travel. They must be weighted graphs.

\* Brute Force Method: The only way to find the optimal solution. It requires that you find the weight of each & every route.

So how many possible routes are there?

3! routes = 6 but some are mirror images of each other.

So  $\frac{6}{2} = 3$  distinct routes



Nearest Neighbors

A W C L H  
3 + 5 + 10 + 7 = 25

Just keep going to the nearest neighbor w/o repetition

Repetitive Nearest Neighbor  
 Do nearest Neighbor Starting from each vertex  
 Find the one of the lowest total weight. Then  
 Rewrite that starting from starting point

For Prim's example:

A W C L H ZS

L W H C L = 23

W H C L = 23

C H W L C = 23  $\rightarrow$  So just rewrite

A W L C H

Cheapest Link

If given graph you may want to write weights in order.

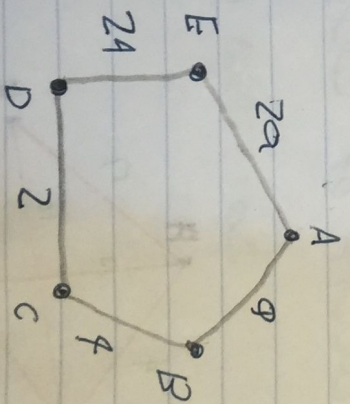
ex.) A B C D E  
 - 9 7 12 29

B 9 - 4 6 14

C 7 4 - 2 11

D 19 6 2 - 19

E 29 14 11 - 24



HW = DO REVIEW